

The graded representation theory of the symmetric group and dominated homomorphisms

Liron Speyer

Queen Mary University of London

l.speyer@qmul.ac.uk

Joint work with Matthew Fayers.

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For each $\lambda \vdash n$, we construct an $\mathbb{F}\mathfrak{S}_n$ -module S_λ , called a Specht module.

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Fact

If $p > 0$ and λ is a *p-regular partition* of n , S_λ has a simple head, which we denote D_λ . These form a complete set of pairwise non-isomorphic simple $\mathbb{F}\mathfrak{S}_n$ -modules.

Graded algebras

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If M is finitely generated, then $\text{Hom}_A(M, N)$ is a graded vector space. In particular, $\text{Hom}_A(M, N)$ has a homogeneous basis.

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$$y_r \psi_r e(i) = (\psi_r y_{r+1} - \delta_{i_r, i_{r+1}}) e(i);$$

$$y_{r+1} \psi_r e(i) = (\psi_r y_r + \delta_{i_r, i_{r+1}}) e(i);$$

$$\psi_r^2 \mathbf{e}(i) = \begin{cases} 0 & i_r = i_{r+1}, \\ \mathbf{e}(i) & i_{r+1} \neq i_r, i_r \pm 1, \\ (y_{r+1} - y_r) \mathbf{e}(i) & i_r = i_{r+1} + 1, \\ (y_r - y_{r+1}) \mathbf{e}(i) & i_r = i_{r+1} - 1, \\ (y_{r+1} - y_r)(y_r - y_{r+1}) \mathbf{e}(i) & i_r = i_{r+1} + 1, p = 2; \end{cases}$$

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$$y_1 = 0;$$

$$\mathbf{e}(i) = 0 \quad \text{if } i_1 \neq 0;$$

for all admissible r, s, i, j .

Cyclotomic KLR algebras and $\mathbb{F}\mathfrak{S}_n$

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If $p > 0$, R_n is non-trivially \mathbb{Z} -graded by setting

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Corollary

If $p > 0$, $\mathbb{F}\mathfrak{S}_n$ can be non-trivially \mathbb{Z} -graded.

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Example: $T_{(3,2)} = \begin{array}{|c|c|c|} \hline 1 & 3 & 5 \\ \hline 2 & 4 & \\ \hline \end{array}$.

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$$\deg(T_{<2}) = \mathbf{0}$$

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We define a degree function $\deg : \text{Std}(\lambda) \rightarrow \mathbb{Z}$ recursively by setting $\deg(T) = \deg^A(\lambda) + \deg(T_{<n})$ where A is the node of T containing n .

Example: Let $p = 3$, $T = \begin{array}{|c|c|c|} \hline 1 & 2 & 5 \\ \hline 3 & 6 & \\ \hline 4 & & \\ \hline \end{array}$. Recursively, $\boxed{0} \boxed{1}$

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$$\deg(T_{<5}) = 0 + 0 + 1 + -1$$

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$$\deg(T) = 0 + 0 + 1 + -1 + 0 + 1 = 1.$$

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For each $w \in \mathfrak{S}_n$, fix a reduced expression $w = s_{i_1} s_{i_2} \dots s_{i_r}$ in the Coxeter generators of \mathfrak{S}_n .

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Example: Take $\lambda = (3, 1)$. We have three standard λ -tableaux

$$\begin{array}{|c|c|c|} \hline 1 & 3 & 4 \\ \hline 2 & & \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline 1 & 2 & 4 \\ \hline 3 & & \\ \hline \end{array} \text{ and } \begin{array}{|c|c|c|} \hline 1 & 2 & 3 \\ \hline 4 & & \\ \hline \end{array}.$$

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1	2	4
3		

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1	2	3
4		

. The corresponding elements v_T are given by z_λ , $\psi_2 z_\lambda$ and $\psi_3 \psi_2 z_\lambda$ respectively.

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Example: $\{z_\lambda, \psi_2 z_\lambda, \psi_3 \psi_2 z_\lambda\}$ is a homogeneous basis of $S_{(3,1)}$.

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2	4		
6			

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residue sequence i_λ . So we check if v_T satisfies the same relations as z_λ (y_r 's all annihilate it, etc). It does, so there's a (degree 1) homomorphism $\mathcal{S}_\lambda \rightarrow \mathcal{S}_\mu$ given by $z_\lambda \mapsto v_T$.

Homomorphisms

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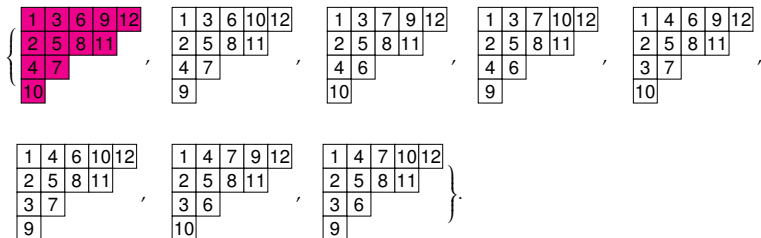
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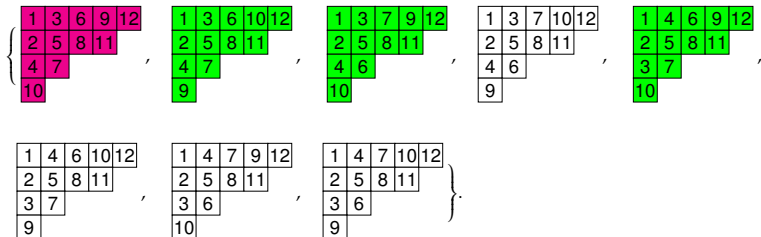
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One of these has degree -2 .

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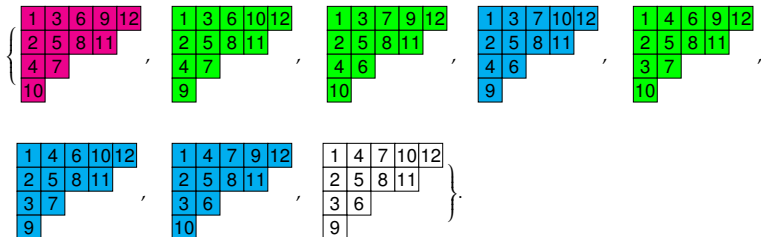
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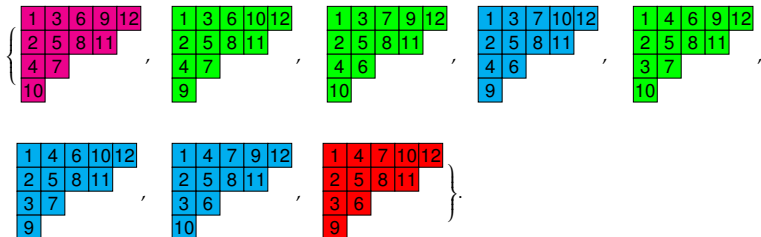
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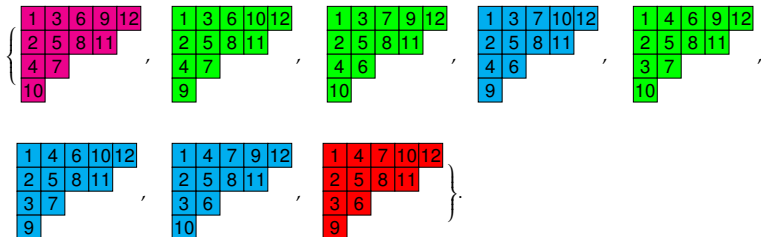
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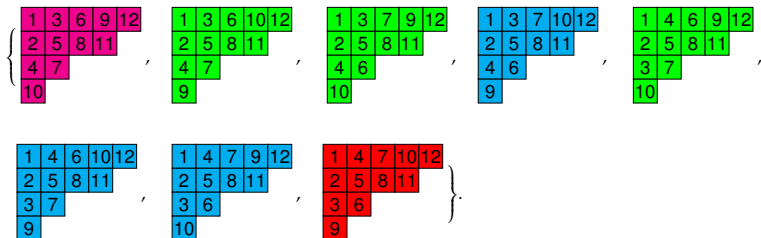


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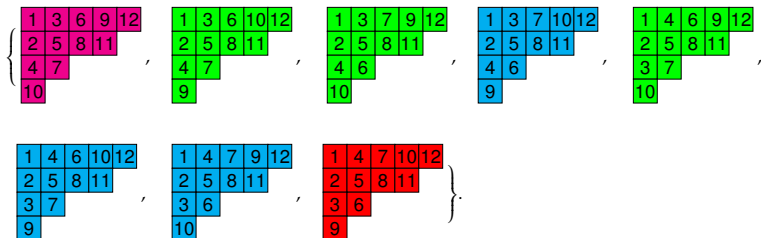


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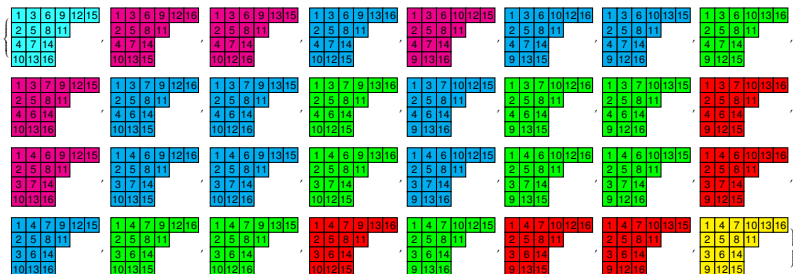
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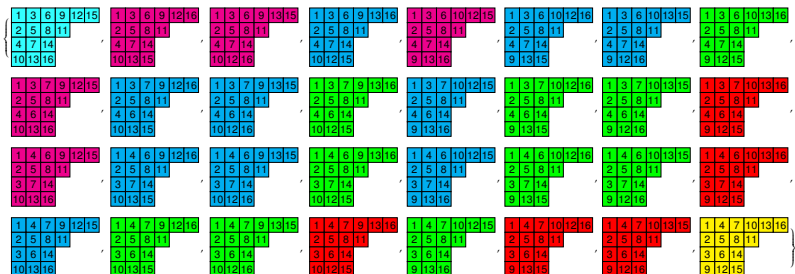
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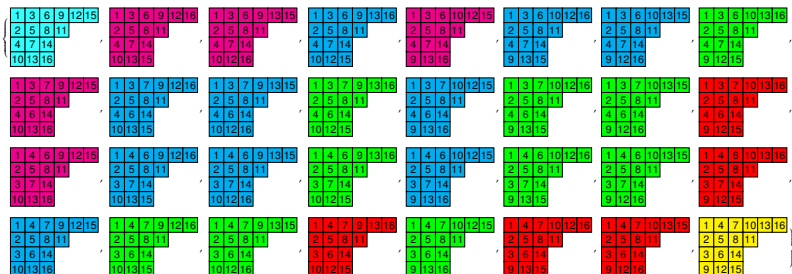
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32 tableaux. Degrees split them into sets of size 1, 1, 5, 5, 10, 10.

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Computation is getting tough...

Dominated tableaux

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Definition

Let $\lambda, \mu \vdash n$. A μ -tableau T is λ -dominated if every entry of T appears at least as far left as it does in T_λ .

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1	3	5	7
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Theorem (Fayers–S, '14)

If $p \neq 2$, then $\text{DHom}_{R_n}(\mathbb{S}_\lambda, \mathbb{S}_\mu) = \text{Hom}_{R_n}(\mathbb{S}_\lambda, \mathbb{S}_\mu)$.

Dominated homomorphisms

Example 2 revisited:

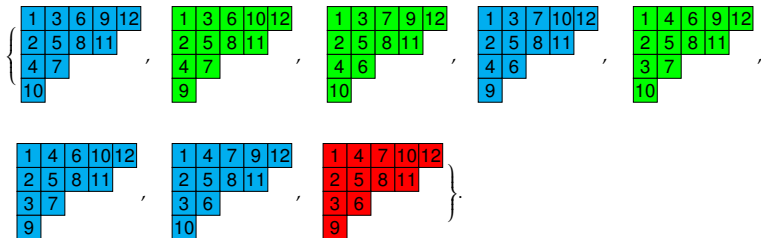
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1	4	7	10	12
2	5	8	11	
3	6			
9				

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Much easier to find the degree 1 homomorphism now!

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Checking the relations is now much easier, and it's not too difficult to show that there is a homomorphism $z_\lambda \mapsto v_T$.

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If $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_k)$, define $\bar{\lambda} := (\lambda_1 - 1, \lambda_2 - 1, \dots, \lambda_k - 1)$. i.e. $\bar{\lambda}$ is the result of removing the first column from λ .

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Theorem (Fayers–S,'14)

If $\lambda, \mu \vdash n$ and λ and μ both have a first column of size k , then

$$\mathrm{DHom}_{R_n}(\mathbf{S}_\lambda, \mathbf{S}_\mu) \cong \mathrm{DHom}_{R_{n-k}}(\mathbf{S}_{\bar{\lambda}}, \mathbf{S}_{\bar{\mu}})$$

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Corollary

If $p \neq 2$, $\lambda, \mu \vdash n$ and λ and μ both have a first column of size k , then

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